

The degree of functions in the Johnson and q -Johnson schemes

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joint work with Jonathan Mannaert and Alfred Wassermann

Introductory remarks

- ▶ Joint work with
[Jonathan Mannaert](#) and [Alfred Wassermann](#).
- ▶ Despite title
“The degree of functions
in the Johnson and q -Johnson [schemes](#)”
No association schemes in this talk!
- ▶ Motivation (next slide) is geometric.
Indeed: Topic close to design theory.
Studied objects are “[dual designs](#)”.

Cameron-Liebler line classes

- ▶ Cameron, Liebler 1982:
“Special” set \mathcal{L} of **lines** in $\text{PG}(3, q)$.
- ▶ Defined by the following equivalent properties:
 - ▶ Algebraic property:
 $\chi_{\mathcal{L}} \in \mathbb{R}$ -row space of the **point-line** incidence matrix.
 - ▶ Geometric property:
Constant intersection with any line spread of $\text{PG}(3, q)$.

In literature: Various directions of generalization

- ▶ Ambient space $\text{PG}(n, q)$.
- ▶ **lines** \longrightarrow k -spaces.
- ▶ Allow $q = 1$ (set case).
- ▶ **points** \longrightarrow spaces of degree t .

Goal

Coherent theory of all above generalizations.

Subset and subspace lattices

- ▶ Fix $q = 1$ (**set case**) or prime power $q \geq 2$ (**q -analog case**).
- ▶ Fix n non-negative integer.
- ▶ Let V be a $\begin{cases} \text{set of size } n \\ \mathbb{F}_q\text{-vector space of dimension } n \end{cases}$
- ▶ Let $\mathcal{L}(V)$ be the lattice of all $\begin{cases} \text{subsets of } V \\ \mathbb{F}_q\text{-subspaces of } V \end{cases}$
- ▶ For $U \in \mathcal{L}(V)$ let $\text{rk}(U) = \begin{cases} \#U \\ \dim(U) \end{cases}$
- ▶ Let $\begin{bmatrix} V \\ k \end{bmatrix} = \{U \in \mathcal{L}(V) \mid \text{rk}(U) = k\}$.
Set case: $\# \begin{bmatrix} V \\ k \end{bmatrix} = \binom{n}{k} = \begin{bmatrix} n \\ k \end{bmatrix}_1$ Binomial coefficient.
 q -analog case: $\# \begin{bmatrix} V \\ k \end{bmatrix} = \begin{bmatrix} n \\ k \end{bmatrix}_q$ Gaussian coefficient.
- ▶ Always: Use **algebraic** dimension!
(Except in established symbols like $\text{PG}(n, q)$).

Algebraic property

- ▶ Algebraic property of Cameron-Liebler line classes:
 $\chi_{\mathcal{L}} \in \mathbb{R}$ -row space of the point-line incidence matrix.
- ▶ Straightforward generalization:
 - ▶ Let $W^{(tk)}$ incidence matrix of t -spaces vs. k -spaces.
 - ▶ Let V_t be the \mathbb{R} -row space of $W^{(tk)}$.
 - ▶ Function $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$ has algebraic property A_t if $f \in V_t$.

Baby example

- ▶ Let $q = 1$, $V = \{1, 2, 3, 4, 5, 6\}$ (so $n = 6$), $k = 3$, $t = 2$.
- ▶ Let $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\} \subseteq \begin{bmatrix} V \\ 3 \end{bmatrix}$.
- ▶ Claim: Set \mathcal{F} has algebraic property A_2 ,
i. e. its characteristic function $\chi_{\mathcal{F}} : \begin{bmatrix} V \\ 3 \end{bmatrix} \rightarrow \mathbb{R}$ has prop. A_2 .

Baby example (cont.)

$$W^{(23)} = \begin{array}{c} 123 \ 124 \ 125 \ 126 \ 134 \ 135 \ 136 \ 145 \ 146 \ 156 \ 234 \ 235 \ 236 \ 245 \ 246 \ 256 \ 345 \ 346 \ 356 \ 456 \\ \left[\begin{array}{cccccccccccccccccccc} 12 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 13 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 14 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 15 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 16 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 23 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 24 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 25 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 \\ 26 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 0 \\ 34 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 35 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 36 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 \\ 45 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 46 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 56 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \end{array} \right] \begin{array}{l} 1/3 \\ 1/3 \\ -1/6 \\ -1/6 \\ -1/6 \\ 1/3 \\ -1/6 \\ -1/6 \\ -1/6 \\ -1/6 \\ -1/6 \\ 1/3 \\ 1/3 \\ 1/3 \end{array} \end{array}$$

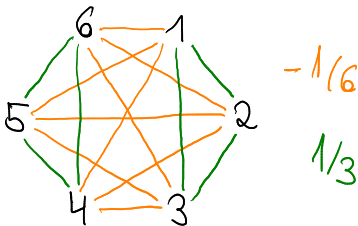
$$\chi_F = (1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1)$$

Baby example (cont.)

- ▶ $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\}$.
- ▶ We found: \mathcal{F} has property A_2 and the vector of **2-weights** of \mathcal{F} is

$$\text{wt}_{\mathcal{F}}^{(2)} = \left(\frac{1}{3}, \frac{1}{3}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, \frac{1}{3}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, \frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right).$$

- ▶ Visualization.



- ▶ Exercise.

\mathcal{F} does **not** have A_1 .

Geometric property

- ▶ Geometric property of Cameron-Liebler line classes:
Constant intersection with any **line** spread of $\text{PG}(3, q)$
- ▶ Generalization? – Not so clear.
- ▶ Observation:
 - line** spread of $\text{PG}(3, q)$
 - = set of **lines** in $\text{PG}(3, q)$ covering every **point** exactly once
 - = simple $1-(4, 2, 1)_q$ subspace design
- ▶ \rightsquigarrow use designs!

Definition: Simple design

A set $\mathcal{D} \subseteq \binom{V}{k}$ is called a **simple t - $(n, k, \lambda)_q$ design**, if every $T \in \binom{V}{t}$ is contained in exactly λ elements of \mathcal{D} .

- ▶ set case $q = 1$: **combinatorial design**
- ▶ q -analog case $q \geq 2$: **subspace design**

Example

▶ Let $q = 1$, $V = \{1, 2, 3, 4, 5, 6\}$ (so $n = 6$), $k = 3$, $t = 2$.

▶ Let

$$\mathcal{D} = \left\{ \{1, 2, 3\}, \{1, 2, 4\}, \{1, 3, 6\}, \{1, 4, 5\}, \{1, 5, 6\}, \right. \\ \left. \{2, 4, 6\}, \{2, 5, 6\}, \{2, 3, 5\}, \{3, 4, 5\}, \{3, 4, 6\} \right\} \subseteq \begin{bmatrix} V \\ 3 \end{bmatrix}.$$

▶ Check design condition for $t = 2$.

▶ $T = \{1, 2\}$ is contained in blocks $\{1, 2, 3\}$ and $\{1, 2, 4\}$.

▶ $T = \{1, 3\}$ is contained in blocks $\{1, 2, 3\}$ and $\{1, 3, 6\}$.

▶ ...

▶ $T = \{5, 6\}$ is contained in blocks $\{1, 5, 6\}$ and $\{2, 5, 6\}$.

▶ $\implies \mathcal{D}$ is simple 2 - $(6, 3, 2)_1$ design.

Example (Trivial simple designs)

▶ \emptyset is empty t - $(v, k, 0)_q$ design.

▶ $\begin{bmatrix} V \\ k \end{bmatrix}$ is complete t - $(v, k, \lambda_{\max})_q$ design
where $\lambda_{\max} := \begin{bmatrix} n-t \\ k-t \end{bmatrix}$.

Definition: Simple design (repeated)

A set $\mathcal{D} \subseteq \binom{V}{k}$ is called a **simple t - $(n, k, \lambda)_q$ design**, if every $T \in \binom{V}{t}$ is contained in exactly λ elements of \mathcal{D} .

- ▶ set case $q = 1$: **combinatorial design**
- ▶ q -analog case $q \geq 2$: **subspace design**

Reformulation in characteristic functions

- ▶ Let \mathbf{x}_T be characteristic function of pencil $\{K \in \binom{V}{k} \mid T \subseteq K\}$.
- ▶ For $f, g : \binom{V}{k} \rightarrow \mathbb{R}$
fix standard inner product $\langle f, g \rangle = \sum_{K \in \binom{V}{k}} f(K)g(K)$.
- ▶ Note that $\#(\mathcal{F} \cap \mathcal{G}) = \langle \chi_{\mathcal{F}}, \chi_{\mathcal{G}} \rangle$ for $\mathcal{F}, \mathcal{G} \subseteq \binom{V}{k}$.
- ▶ \mathcal{D} is simple t - $(n, k, \lambda)_q$ design
 $\iff \langle \mathbf{x}_T, \chi_{\mathcal{D}} \rangle = \lambda$ for all $T \in \binom{V}{t}$.
- ▶ \rightsquigarrow generalization to **real designs**.

Generalized definition: Real design

A function $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$ is called a **real t - $(n, k, \lambda)_q$ design**, if $\langle \mathbf{x}_T, f \rangle = \lambda$ for all $T \in \begin{bmatrix} V \\ t \end{bmatrix}$.

- ▶ **f null design or trade** if $\lambda = 0$.
- ▶ **f signed design** if $\text{im}(f) \subseteq \mathbb{Z}$.
- ▶ **f design or possibly non-simple design** if $\text{im}(f) \subseteq \mathbb{N}$.
(Idea: simple design, but with possibly repeated blocks)
- ▶ **f (characteristic function of) simple design**
 $\iff \text{im}(f) \subseteq \{0, 1\} \iff f$ **Boolean**.

Further reformulation

- ▶ **Observation:**
Functions \mathbf{x}_T (interpreted as vectors)
are the rows of incidence matrix $W^{(tk)}$.
- ▶ **Therefore:**
 f real t - $(n, k, \lambda)_q$ design $\iff W^{(tk)}f = \lambda \mathbf{1}$.
- ▶ **In particular:**
 f real t - $(n, k, 0)_q$ null design $\iff W^{(tk)}f = \mathbf{0}$.

Geometric property, basic version

- ▶ For $\lambda \in \mathbb{R}$ let $U_\lambda :=$ set of real t - $(n, k, \lambda)_q$ design.
- ▶ Just seen: $U_0 = \ker W^{(tk)}$.
- ▶ Set of functions with A_t was $V_t = \text{rowsp } W^{(tk)}$.

$$\implies V_t = U_0^\perp$$

What did we get?

- ▶ Established a connection to designs.
- ▶ Concept known as Delsarte's **design orthogonality**.
- ▶ Compared to prototype
“constant intersection with all spreads”:
Want similar property for $\lambda \neq 0$!

Geometric property, version II

- ▶ Fix $\lambda \in \mathbb{R}$.
- ▶ Scaled complete design $\frac{\lambda}{\lambda_{\max}} \cdot \mathbf{1}$ is real t -(n, k, λ) $_q$ design.
- ▶ As solution of linear equation system $W^{(tk)}f = \lambda \mathbf{1}$:

$$U_\lambda = \frac{\lambda}{\lambda_{\max}} \cdot \mathbf{1} + \underbrace{\ker W^{(tk)}}_{=U_0=V_t^\perp}.$$

- ▶ \implies

$$U_\lambda = \left\{ \delta : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R} \mid \langle f, \delta \rangle = \frac{\lambda}{\lambda_{\max}} \cdot \#f \text{ for all } f \in V_t \right\} \quad \text{and}$$

$$V_t = \left\{ f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R} \mid \langle f, \delta \rangle = \frac{\lambda}{\lambda_{\max}} \cdot \#f \text{ for all } \delta \in U_\lambda \right\} \quad \text{Vers. II}$$

(with $\#f = \sum_{K \in \begin{bmatrix} V \\ k \end{bmatrix}} f(K) = \langle f, \mathbf{1} \rangle$, motivated by $\#\mathcal{F} = \#\chi_{\mathcal{F}}$)

- ▶ Still room for improvement:
 - ▶ Not happy about “For all **real** ... designs”.
 \rightsquigarrow enough to look at **basis** of U_λ .
 - ▶ Allow mixed values of λ .

Example

- ▶ $q = 1, n = 6, k = 3, t = 2 \rightsquigarrow \lambda_{\max} = \begin{bmatrix} 6-2 \\ 3-2 \end{bmatrix} = 4$.
- ▶ Baby example: $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\}$, seen: $\chi_{\mathcal{F}} \in V_2$.
- ▶ Geometric property \implies For each 2 -(6, 3, 2)₁ design:

$$\langle \chi_{\mathcal{F}}, \delta \rangle = \frac{\lambda}{\lambda_{\max}} \cdot \#\chi_{\mathcal{F}} = \frac{2}{4} \cdot 2 = 1.$$

- ▶ \implies Each **simple** 2 -(6, 3, 2)₁ design \mathcal{D} contains **exactly one** of the blocks $\{1, 2, 3\}$ and $\{4, 5, 6\}$.
- ▶ $\rightsquigarrow \mathcal{D}$ is **anti-complementary**.
- ▶ Can also be shown using **intersection numbers**.

Geometric property, toolbox version

- ▶ U_* := set of all real t - $(v, k, \lambda)_q$ designs with arbitrary value $\lambda \in \mathbb{R}$.
- ▶ By scaled complete designs: $U_* = U_0 + \langle \mathbf{1} \rangle_{\mathbb{R}}$.
- ▶ Lemma (Toolbox version of geometric property).
Let $\Delta \subseteq U_*$. Then

$$V_t = \left\{ f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R} \mid \langle f, \delta \rangle = \frac{\lambda_\delta}{\lambda_{\max}} \cdot \#f \text{ for all } \delta \in \Delta \right\}$$
$$\iff \langle \Delta \cup \{ \mathbf{1} \} \rangle_{\mathbb{R}} = U_*$$

Proof. Dimension argument. Use that $W^{(tk)}$ has full rank.
(Set case: Gottlieb 1966, q -analog case: Kantor 1972)

- ▶ Question: Suitable sets Δ ?

Lemma

Let Δ be

- (a) the set of all signed t - $(n, k, 0)_q$ null designs or
- (b) the set of all possibly non-simple t - $(n, k, \lambda)_q$ designs

Then $U_* = \langle \Delta \cup \{\mathbf{1}\} \rangle_{\mathbb{R}}$.

Proof.

Part (a).

- ▶ entries of $W^{(tk)}$ are in \mathbb{Q} .
- ▶ $\implies U_0 = \ker W^{(tk)}$ has rational basis.
- ▶ Multiply by common denominators \rightsquigarrow integral basis B .
- ▶ $\implies B \subseteq \Delta$ and $\langle B \cup \{\mathbf{1}\} \rangle_{\mathbb{R}} = U_*$.

Part (b).

- ▶ Start with B .
- ▶ Add suitable integral multiples of $\mathbf{1}$
 \rightsquigarrow non-negative integral set B' .
- ▶ $\implies B' \subseteq \Delta$ and $\langle B' \cup \{\mathbf{1}\} \rangle_{\mathbb{R}} = U_*$.

We arrive at:

Theorem

Let $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$. The following are equivalent.

(i) *Algebraic property:* $f \in V_t$.

Geometric properties:

(ii) There is a constant $c \in \mathbb{R}$ such that $\langle f, \delta \rangle = \lambda_\delta c$
for all *real* t - $(n, k, \lambda_\delta)_q$ designs δ with $\lambda_\delta \in \mathbb{R}$.

(iii) $\langle f, \delta \rangle = 0$
for all *signed* t - $(n, k, 0)_q$ *null* designs $\delta : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{Z}$.

(iv) There is a constant $c \in \mathbb{R}$ such that $\langle f, \delta \rangle = \lambda_\delta c$
for all *possibly non-simple* t - $(n, k, \lambda_\delta)_q$ designs $\delta : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{N}$.

The constant in properties (ii) and (iv) necessarily equals

$$c = \frac{1}{\lambda_{\max}} \cdot \#f.$$

Geometric property: Discussion

- ▶ Tempting: Is the following a suitable geometric property?

“There is a constant $c \in \mathbb{R}$ such that $\langle f, \delta \rangle = \lambda c$ for all **simple** t - $(n, k, *)_q$ designs”

- ▶ By toolbox version: If and only if $\langle \{\text{simple } t\text{-}(n, k, *)_q \text{ designs}\} \rangle_{\mathbb{R}} = U_*$ (**richness cond**)
- ▶ Unfortunately: Not always true.

Counterexample. $q = 1, n = 10, k = 5, t = 4$.

By integrality conditions: All simple 4 - $(10, 5, *)_1$ are trivial.

$\implies \dim \langle \{\text{simple } 4\text{-}(10, 5, *)_1 \text{ designs}\} \rangle_{\mathbb{R}} = 1$, too small!

- ▶ **Research problem.** (probably hard!)

Classify the parameters (q, n, k, t) where the richness condition holds.

The Degree

▶ Fix $k \in \{0, \dots, n\}$ and $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$.

▶ Lemma.

$$\{\mathbf{1}\} = V_0 \subsetneq V_1 \subsetneq \dots \subsetneq V_k = V.$$

Proof. $W^{(ij)}W^{(jk)} \sim W^{(ik)}$ for $0 \leq i \leq j \leq k$.

▶ Definition.

Degree $\deg(f) :=$ smallest t such that $f \in V_t$.

Example

▶ Functions f of degree 0

are the **scalar functions** $f = \lambda \mathbf{1}$ with $\lambda \in \mathbb{R}$.

▶ **Baby example** $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\}$.

In $V = \{1, 2, 3, 4, 5, 6\}$ we have $\deg(\mathcal{F}) := \deg(\chi_{\mathcal{F}}) = 2$.

▶ Seen: $\chi_{\mathcal{F}} \in V_2$.

▶ Exercise: $\chi_{\mathcal{F}} \notin V_1$.

▶ In $V = \{1, 2, 3, 4, 5, 6, 7\}$ we have $\deg(\mathcal{F}) = 3$.

\implies **Ambient space** V matters!

The Degree (cont.)

- ▶ **Remember.** Rows of $W^{(tk)}$ are the t -pencils \mathbf{x}_T .
- ▶ \rightsquigarrow **Alternative characterization of degree.**

$\deg(f)$ is smallest t

such that f is a linear combination of t -pencils \mathbf{x}_T .

The (unique) coefficients are called **t -weights** $\text{wt}_f(T)$ of f :

$$f = \sum_{T \in \binom{V}{t}} \text{wt}_f(T) \mathbf{x}_T$$

Lemma

- (a) $\deg(\lambda f) \leq \deg(f)$ *with equality iff* $\lambda \neq 0$.
- (b) $\deg(f + g) \leq \max(\deg(f), \deg(g))$.
- (c) $\deg(fg) \leq \deg(f) + \deg(g)$.

Proof.

Parts (a), (b): easy. Part (c): Use weights & $\deg \mathbf{x}_T \leq \text{rk } T$. □

Dualization

- ▶ Fix anti-isomorphism \perp of the lattice $\mathcal{L}(V)$.
 - ▶ Set case: Set complement.
 - ▶ q -analog case: Perp wrt non-degenerate bilinear form.
- ▶ Induces **dual map** of $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$:

$$f^\perp : \begin{bmatrix} V \\ n - k \end{bmatrix} \rightarrow \mathbb{R}, \quad U \mapsto f(U^\perp)$$

- ▶ Effect of dualization on the degree?

Theorem

(a) $\deg f^\perp = \deg f$.

(b) For $i \in \{0, \dots, \deg f\}$, the i -weight distribution of f^\perp is

$$\text{wt}_{f^\perp}^{(i)}(\mathcal{J}) = \sum_{I \in \binom{V}{i}} \gamma(n-k, i, \text{rk}(I^\perp \cap \mathcal{J})) \text{wt}_f^{(i)}(I)$$

where

$$\gamma(k, i, z) := \begin{cases} \delta_{z,k} & \text{if } i = k, \\ (-1)^{i-z} \frac{1}{q^{(k-i)(i-z) + \binom{i-z}{2}}} \frac{\begin{bmatrix} k-i \\ 1 \end{bmatrix}}{\begin{bmatrix} k-z \\ 1 \end{bmatrix}} \frac{1}{\begin{bmatrix} k \\ z \end{bmatrix}} & \text{otherw.} \end{cases}$$

Proof.

- ▶ Enough to look at pencils $f = \mathbf{x}_J$.
- ▶ Set up linear equation system for the weights of f^\perp , assuming that $\text{wt}(I)$ only depends on $\text{rk}(I \cap \mathcal{J})$.
- ▶ Equation system matrix is triangular with non-zero diagonal \implies invertible \implies Part (a).
- ▶ Apply negation formula & q -Vandermonde formula for Gaussian coefficients \rightsquigarrow compute solution \rightsquigarrow Part (b).

Change of ambient space

Two elementary ways to shrink the ambient space V .

- ▶ $V \rightarrow H$ ($H \in \begin{bmatrix} V \\ n-1 \end{bmatrix}$ hyperplane)
- ▶ $V \rightarrow V/P$ ($P \in \begin{bmatrix} V \\ 1 \end{bmatrix}$ point)

Implication on the degree?

We start with $V \rightarrow V/P$.

Theorem

Let $1 \leq k \leq n$ and $P \in \begin{bmatrix} V \\ 1 \end{bmatrix}$. Then

$$\Phi : \mathbb{R}^{\begin{bmatrix} V/P \\ k-1 \end{bmatrix}} \rightarrow \mathbb{R}^{\begin{bmatrix} V \\ k \end{bmatrix}}, \quad \Phi(f) : K \mapsto \begin{cases} f(K/P) & \text{if } P \subseteq K, \\ 0 & \text{if } P \not\subseteq K \end{cases}$$

is an injective \mathbb{R} -linear map with

$$\text{im}(\Phi) = \left\{ g \in \mathbb{R}^{\begin{bmatrix} V \\ k \end{bmatrix}} \mid \text{supp } g \subseteq \begin{bmatrix} V \\ k \end{bmatrix} \Big|_P \right\} \quad \text{and}$$

$$\deg_V \Phi(f) = \begin{cases} 0 & \text{if } f = 0, \\ \min(\overbrace{\deg_{V/P}(f) + 1}^{\text{main case}}, n - k) & \text{otherwise.} \end{cases}$$

Proof.

- ▶ Straightforward, except “ $\deg_V \Phi(f) \geq \deg_{V/P}(f) + 1$ ”.
- ▶ **Lemma.** In main case

For all $g \in \text{im } \Phi$: $P \not\subseteq T \implies \text{wt}_g(T) = 0$.

Proof. Incidence matrices of certain attenuated geometries are of full rank. (Guo, Li, Wang, 2014.)

Theorem

Let $1 \leq n - k \leq n$ and $H \in \binom{V}{n-1}$. Then

$$\psi : \mathbb{R}^{\binom{H}{k}} \rightarrow \mathbb{R}^{\binom{V}{k}}, \quad \psi(f) : K \mapsto \begin{cases} f(K) & \text{if } K \subseteq H, \\ 0 & \text{if } K \not\subseteq H \end{cases}$$

is an injective \mathbb{R} -linear map with

$$\begin{aligned} \text{im}(\psi) &= \{g \in \mathbb{R}^{\binom{V}{k}} \mid \text{supp } g \subseteq \binom{H}{k}\} \quad \text{and} \\ \text{deg}_V \psi(f) &= \begin{cases} 0 & \text{if } f = 0, \\ \min(\text{deg}_H(f) + 1, k) & \text{otherwise.} \end{cases} \end{aligned}$$

Proof.

Follows from the previous theorem by [dualization](#). □

Example (Basic sets)

- ▶ Start with “complete set” $\begin{bmatrix} W \\ \ell \end{bmatrix}$ of degree 0.



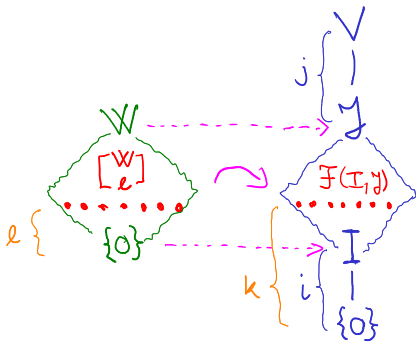
- ▶ i -fold application of Φ and j -fold application of Ψ

$$\rightsquigarrow \text{basic set } \mathcal{F}(I, J) = \left\{ K \in \begin{bmatrix} V \\ k \end{bmatrix} \mid I \subseteq K \subseteq J \right\}.$$

- ▶ By theorems: $\deg \mathcal{F}(I, J) = i + j$ (in the main case).

Example (Basic sets)

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Example (Basic sets (cont.))

- ▶ Basic sets $\mathcal{F}(I, J)$ include pencils ($j = 0$) and dual pencils ($i = 0$).
In particular $\deg \mathbf{x}_I = k$ (in the main case).
- ▶ Geometric property of $\mathcal{F}(I, J)$



design property of i -fold derived and j -fold residual design.

Sets and Boolean functions

- ▶ Of particular interest: **Sets** $\mathcal{F} \subseteq \binom{V}{k}$ of low degree.
- ▶ Via characteristic functions:
Sets correspond to **Boolean** functions $\binom{V}{k} \rightarrow \{0, 1\}$.

Boolean degree 1 functions

- ▶ **Set** case:

Filmus, Ihringer 2019:

Only **basic** functions.

\implies only pencils and dual pencils (since $t = 1$).

- ▶ **q -analog** case:

Boolean degree 1 function = **Cameron-Liebler** set of $(k - 1)$ -spaces in $\text{PG}(n - 1, q)$.

Non-basic examples do exist.

Classification: Hard research problem.

Computer classification

Goal.

For $q = 1$ and small n, k , **classify** all sets \mathcal{F} of degree $t = 2$.

Strategy.

- ▶ Use “basic” **geometric property**:

$$\deg \chi_{\mathcal{F}} \leq t \iff \chi_{\mathcal{F}} \in \ker W^{(tk)}.$$

\rightsquigarrow Want to find all $\{0, 1\}$ -vectors in $\ker W^{(tk)}$.

- ▶ Find **integral basis** of $\ker W^{(tk)}$.
 - ▶ either: computationally
 - ▶ or: Use literature like
Khosrovshahi, Ajoodani-Namini (1990):
A new basis for trades
- ▶ \rightsquigarrow system of linear Diophantine equations.
- ▶ Solve using **SOLVEDIOPHANT** (A. Wassermann)
- ▶ Filter out isomorphic copies.
(action of symmetric group \mathfrak{S}_n)

Results

n	k	size distribution	Σ
6	3	2 4 ³ 6 ⁵ 8 ⁸ 10 ⁸ 12 ⁸ 14 ⁵ 16 ³ 18	42
7	3	5 ² 10 ⁶ 15 ¹¹ 20 ¹¹ 25 ⁶ 30 ²	38
8	3	6 8 11 12 14 15 ² 16 17 18 20 ² 21 22 23 24 ³ 25 26 ⁴ 27 28 ² ...	50
9	3	7 14 ² 21 ⁵ 28 ⁴ 35 ⁵ 42 ¹¹ 49 ⁵ 56 ⁴ 63 ⁵ 70 ² 77	45
10	3	8 16 20 24 ² 28 ³ 32 ² 36 ³ 40 ² 44 ² 48 ² 52 ² 56 ⁵ 60 ⁵ 64 ⁵ 68 ² 72 ² ...	57
8	4	10 15 ² 20 ³ 30 ⁶ 35 ² 40 ⁶ 50 ³ 55 ² 60	26
9	4	21 ² 35 ³ 56 ⁴ 70 ⁴ 91 ³ 105 ²	18
10	4	28 42 56 ² 70 84 98 ³ 112 ³ 126 140 154 ² 168 182	18
11	4	36 78 84 ² 120 126 162 ³ 168 ³ 204 210 246 ² 252 294	18
12	4	45 120 ² 135 165 210 240 ³ 255 ³ 285 330 360 375 ² 450	18

blue = sizes of basic sets

Goal. Explain divisibility pattern of the sizes!

Theorem (Divisibility theorem)

Let $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{Z}$ be a function of degree t . Then

$$\underbrace{\gcd \left(\begin{bmatrix} n-0 \\ k-0 \end{bmatrix}, \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}, \dots, \begin{bmatrix} n-t \\ k-t \end{bmatrix} \right)}_{=: a} \mid \#f.$$

Proof.

▶ Algebraic property $\Rightarrow \exists \mathbf{x} : \begin{bmatrix} V \\ t \end{bmatrix} \rightarrow \mathbb{R}$ with $\mathbf{x}^\top W^{(tk)} = f^\top$. (1)

▶ Complete design: $W^{(tk)} \cdot \mathbf{1} = \lambda_{\max} \cdot \mathbf{1}$ (2)

▶ Design theory:

parameters t - $(n, k, \lambda_{\min})_q$ with $\lambda_{\min} = \frac{\lambda_{\max}}{a}$ are **admissible**.

▶ $\Rightarrow \exists$ **signed** t - $(n, k, \lambda_{\min})_q$ design $\delta \Rightarrow W^{(tk)}\delta = \lambda_{\min} \cdot \mathbf{1}$ (3)

▶ Set case: Wilson, "The necessary conditions for t -designs are sufficient for something" (1973).

▶ q -analog case: Ray-Chaudhuri, Singhi (1989).

▶ Left multiplication of (2) and (3) by \mathbf{x}^\top , using (1)

$$\implies \#f = \lambda_{\max} \cdot \#\mathbf{x} \quad \text{and} \quad \langle f, \delta \rangle = \lambda_{\min} \cdot \#\mathbf{x}$$

▶ $\implies \#f = a \cdot \underbrace{\langle f, \delta \rangle}_{\in \mathbb{Z}} \in \mathbb{Z}.$

Compare with the results

$$q = 1, t = 2 \implies a = \gcd\left(\binom{n}{k}, \binom{n-1}{k-1}, \binom{n-2}{k-2}\right).$$

n	k	size distribution	a
6	3	$2 \cdot 4^3 \cdot 6^5 \cdot 8^8 \cdot 10^8 \cdot 12^8 \cdot 14^5 \cdot 16^3 \cdot 18$	2
7	3	$5^2 \cdot 10^6 \cdot 15^{11} \cdot 20^{11} \cdot 25^6 \cdot 30^2$	5
8	3	$6 \cdot 8 \cdot 11 \cdot 12 \cdot 14 \cdot 15^2 \cdot 16 \cdot 17 \cdot 18 \cdot 20^2 \cdot 21 \cdot 22 \cdot 23 \cdot 24^3 \cdot 25 \cdot 26^4 \cdot 27 \cdot 28^2 \dots$	1
9	3	$7 \cdot 14^2 \cdot 21^5 \cdot 28^4 \cdot 35^5 \cdot 42^{11} \cdot 49^5 \cdot 56^4 \cdot 63^5 \cdot 70^2 \cdot 77$	7
10	3	$8 \cdot 16 \cdot 20 \cdot 24^2 \cdot 28^3 \cdot 32^2 \cdot 36^3 \cdot 40^2 \cdot 44^2 \cdot 48^2 \cdot 52^2 \cdot 56^5 \cdot 60^5 \cdot 64^5 \cdot 68^2 \cdot 72^2 \dots$	4
8	4	$10 \cdot 15^2 \cdot 20^3 \cdot 30^6 \cdot 35^2 \cdot 40^6 \cdot 50^3 \cdot 55^2 \cdot 60$	5
9	4	$21^2 \cdot 35^3 \cdot 56^4 \cdot 70^4 \cdot 91^3 \cdot 105^2$	7
10	4	$28 \cdot 42 \cdot 56^2 \cdot 70 \cdot 84 \cdot 98^3 \cdot 112^3 \cdot 126 \cdot 140 \cdot 154^2 \cdot 168 \cdot 182$	14
11	4	$36 \cdot 78 \cdot 84^2 \cdot 120 \cdot 126 \cdot 162^3 \cdot 168^3 \cdot 204 \cdot 210 \cdot 246^2 \cdot 252 \cdot 294$	6
12	4	$45 \cdot 120^2 \cdot 135 \cdot 165 \cdot 210 \cdot 240^3 \cdot 255^3 \cdot 285 \cdot 330 \cdot 360 \cdot 375^2 \cdot 450$	15

Perfect fit!

Parameter of Cameron-Liebler sets of k -spaces

- ▶ Consider q -analog case $q \geq 2$.
- ▶ For sets \mathcal{F} of degree $t = 1$ define
parameter $x := \#\mathcal{F} / \begin{bmatrix} n-1 \\ k-1 \end{bmatrix} \in \mathbb{Q}$
- ▶ Corollary of divisibility theorem:

$$\frac{q^k - 1}{q^{\gcd(n,k)} - 1} \cdot x \in \mathbb{Z},$$

restricting denominator of fraction x in canceled form.

Example

- ▶ $k \mid n \implies x \in \mathbb{Z}$.
Already known: Blokhuis, De Boeck, D'haeseleer (2019).
- ▶ n and k coprime $\implies (1 + q + \dots + q^{k-1}) \cdot x \in \mathbb{Z}$.
- ▶ $k = 4, n \equiv 2 \pmod{4} \implies (1 + q^2) \cdot x \in \mathbb{Z}$.

The paired construction

- ▶ Construction for the **set case** $q = 1$ only.
- ▶ **Idea.** Disjoint union of two “opposite” basic sets.
- ▶ Let $I, J \subseteq V$ be disjoint, not both empty. Define

$$\mathcal{P}(I, J) := \mathcal{F}(I, J^c) \uplus \mathcal{F}(J, I^c)$$

- ▶ Clear:

$$\deg \mathcal{P}(I, J) \leq \min(\#I + \#J, k) \quad (\text{trivial bound})$$

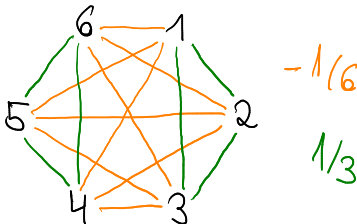
- ▶ Will see: There are cases with a **strict** “ $<$ ”!

Example (1)

- ▶ $V = \{1, \dots, 6\}$, $k = 3$, $I = \emptyset$, $J = \{4, 5, 6\}$,

$$\begin{aligned} & \mathcal{P}(\emptyset, \{4, 5, 6\}) \\ &= \mathcal{F}(\emptyset, \{1, 2, 3\}) \uplus \mathcal{F}(\{4, 5, 6\}, \{1, 2, 3, 4, 5, 6\}) \\ &= \{\{1, 2, 3\}, \{4, 5, 6\}\} \end{aligned}$$

- ▶ This is the **Baby example**!
- ▶ Already seen: $\deg \mathcal{P}(\emptyset, \{1, 2, 3\}) = 2$



- ▶ ... beating the trivial bound " ≤ 3 "!

Example (2)

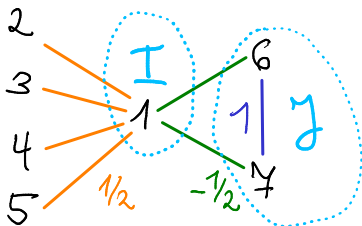
- ▶ $V = \{1, \dots, 7\}$, $k = 3$, $I = \{1\}$, $J = \{6, 7\}$.

$$\mathcal{P}(\{1\}, \{6, 7\})$$

$$= \mathcal{F}(\{1\}, \{1, 2, 3, 4, 5\}) \uplus \mathcal{F}(\{6, 7\}, \{2, 3, 4, 5, 6, 7\})$$

$$= \{\{1, 2, 3\}, \{1, 2, 4\}, \{1, 2, 5\}, \{1, 3, 4\}, \{1, 3, 5\}, \{1, 4, 5\}, \\ \{2, 6, 7\}, \{3, 6, 7\}, \{4, 6, 7\}, \{5, 6, 7\}\}$$

- ▶ Visualization of 2-weights:



- ▶ $\implies \deg(\mathcal{P}(\{1\}, \{6, 7\})) = 2$.
- ▶ ... again beating the trivial bound " ≤ 3 "!

Theorem

Let $q = 1$, $I, J \subseteq V$ disjoint, $i = \#I$, $j = \#J$, $k \leq \frac{n}{2}$,
 $i \leq k \leq n - i$, $j \leq k \leq n - j$.

In the cases

- (a) $i + j \leq k$ and $i + j$ odd;
- (b) $i + j \geq k$ and k odd and $n = 2k$

we have

$$\deg \mathcal{P}(I, J) \leq \min(i + j, k) - 1.$$

Proof (Idea).

Part (a): Write $\chi_{\mathcal{P}(I, J)}$ as an integer linear combination of basic functions of degree $i + j - 1$.

Part (b):

- ▶ Use $\mathcal{P}(X, Y) = \mathcal{P}(X \uplus \{x\}, Y) \uplus \mathcal{P}(X, Y \uplus \{x\})$
(where $X, Y, \{x\}$ are pairwise disjoint)
- ▶ Moving elements from J to $I \rightsquigarrow \deg \mathcal{P}(I, J) \leq \deg \mathcal{P}(K, J')$
- ▶ $\mathcal{P}(K, J') = \mathcal{P}(K, \emptyset) \implies$ Back in Case (a).

Theorem

Let $q = 1$, $I, J \subseteq V$ disjoint, $i = \#I$, $j = \#J$, $k \leq \frac{n}{2}$,
 $i \leq k \leq n - i$, $j \leq k \leq n - j$.

In the cases

- (a) $i + j \leq k$ and $i + j$ odd;
- (b) $i + j \geq k$ and k odd and $n = 2k$

we have

$$\deg \mathcal{P}(I, J) \leq \min(i + j, k) - 1.$$

Work in progress / Conjecture

Statement of Theorem is **best possible**.

- ▶ In fact always equality

$$\deg \mathcal{P}(I, J) = \min(i + j, k) - 1.$$

- ▶ In all cases not covered by (a) and (b),
the trivial bound is sharp:

$$\deg \mathcal{P}(I, J) = \min(i + j, k).$$

Small sets of degree t

▶ Natural question.

Smallest size $m_q(n, k, t)$ of a non-empty set of degree $\leq t$?

▶ From $\deg \mathbf{x}_T = t$ we get

$$m_q(n, k, t) \leq \binom{n-t}{k-t}. \quad (*)$$

▶ Bound (*) is always sharp for $t = 1$.

▶ Set case: Filmus, Ihringer (2019).

▶ q -analog case: Blokhuis, De Boeck, D'haeseleer (2019).

▶ For $q = 1$, $n = 2k$, $t \geq 2$ even, $i = 0$ and $j = t + 1$,
the paired construction beats bound (*)!

Corollary

Let $t \in \{0, \dots, k-1\}$ be even. Then

$$m_1(2k, k, t) \leq 2 \cdot \binom{2k-t-1}{k}.$$

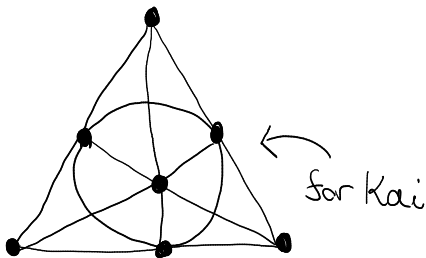
Open problems

- ▶ Many!
- ▶ For fixed (q, n, k, t) , characterize the **sizes** of degree t sets.
 - ▶ Smallest,
 - ▶ second smallest,
 - ▶ gaps,
 - ▶ etc.
- ▶ Further investigate and exploit **relationship**
degree t functions \longleftrightarrow t -designs.

Which results can be translated?

- ▶ Maybe most important:
Better **name** for the studied objects.
 - ▶ “dual designs”? \longrightarrow ambiguous.
 - ▶ Something involving “Cameron-Liebler”?
 - ▶ other ideas?

Thank you!



Slides will be uploaded at

<https://mathe2.uni-bayreuth.de/michaelk/>